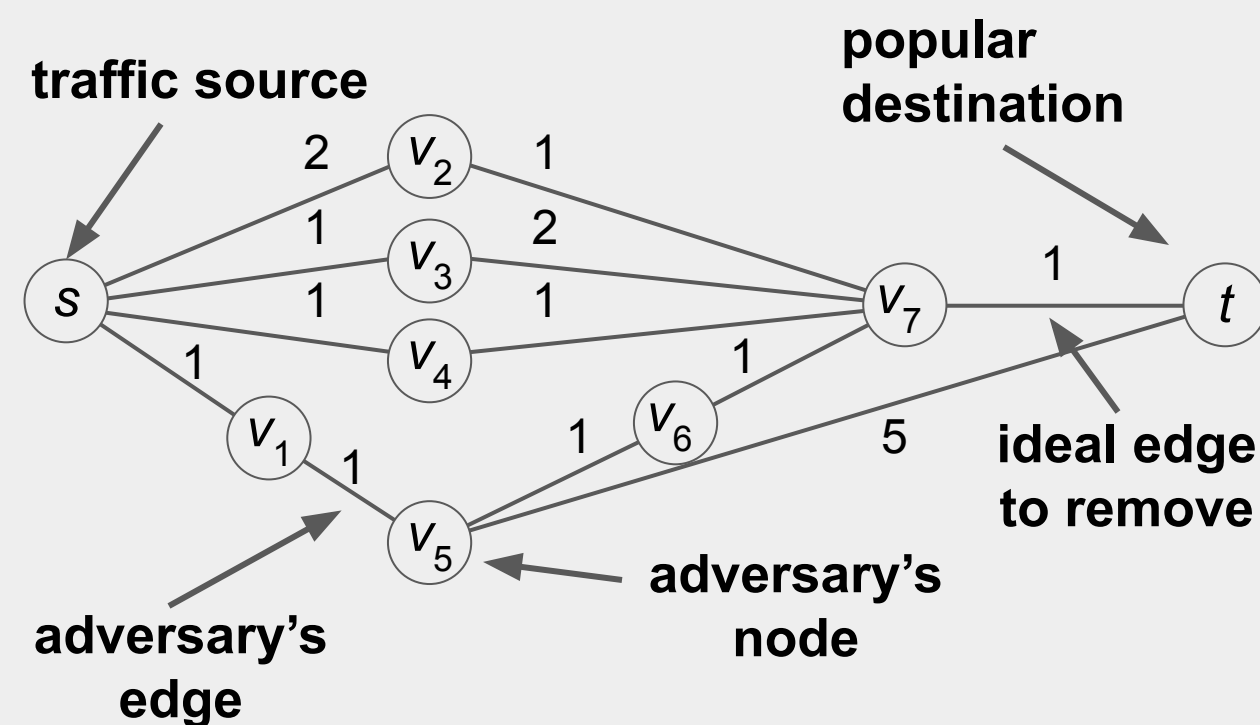


## Problem Statement

Scenario: An adversary can divert traffic to specific parts of a graph by removing edges



**Goal:** Increase attacker's budget required to successfully attack

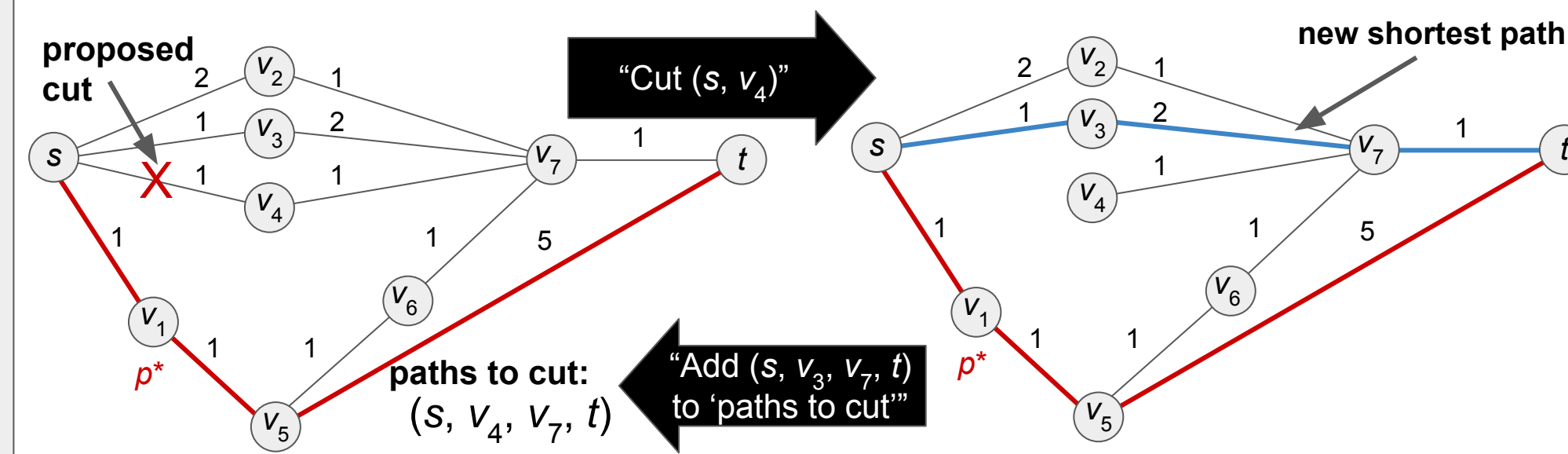
### Assumed Context:

- We have a graph and know the true weights
- But we publish *approximate* weights
- Desired outcome: Adversary's required budget is increased
- Constraint: Minimize negative impact on legitimate users

### Research Questions:

- What is the impact on legitimate network users when we stop an attack?
- Is it more costly to defend against an attack that targets a path that exits and returns to a community?

## Attack: PATHATTACK\*

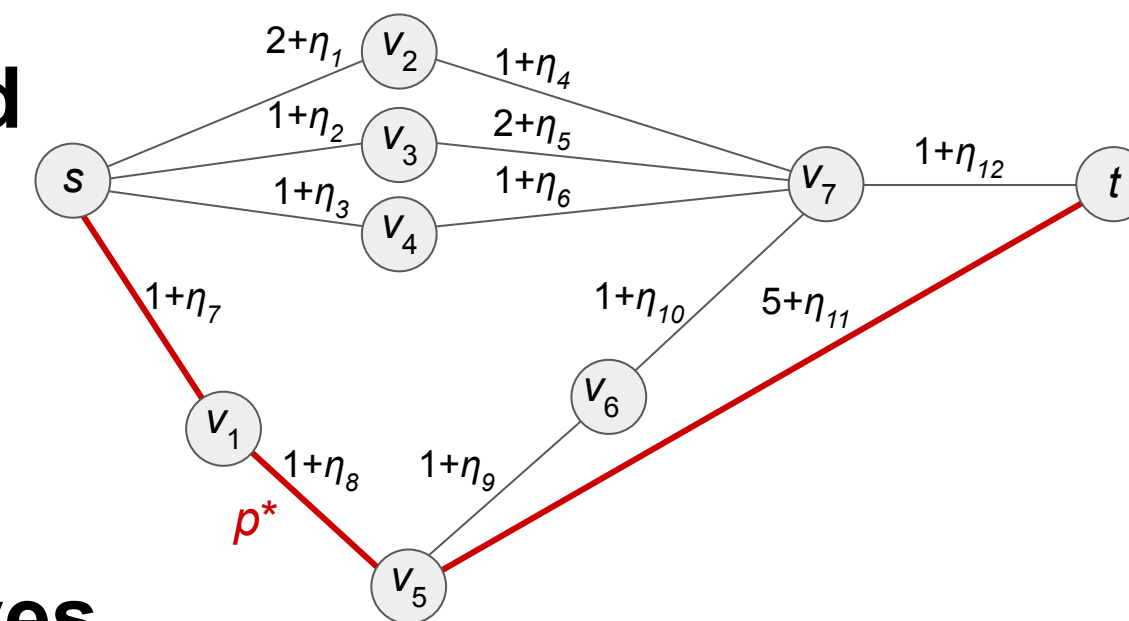


- Algorithm to approximate optimal path cut cost
- Iteratively add constraints until  $p^*$  is shortest

\*B. A. Miller et al., "PATHATTACK: Attacking Shortest Paths in Complex Networks," in ECML PKDD, 2021

## Defense: Modified Weights

- Prior work<sup>†</sup> used noisy weights to increase adversarial uncertainty
- This work involves optimizing the defender's cost function with modified weights



<sup>†</sup>B. A. Miller et al., "Defense Against Shortest Path Attacks," GraphEx, 2022

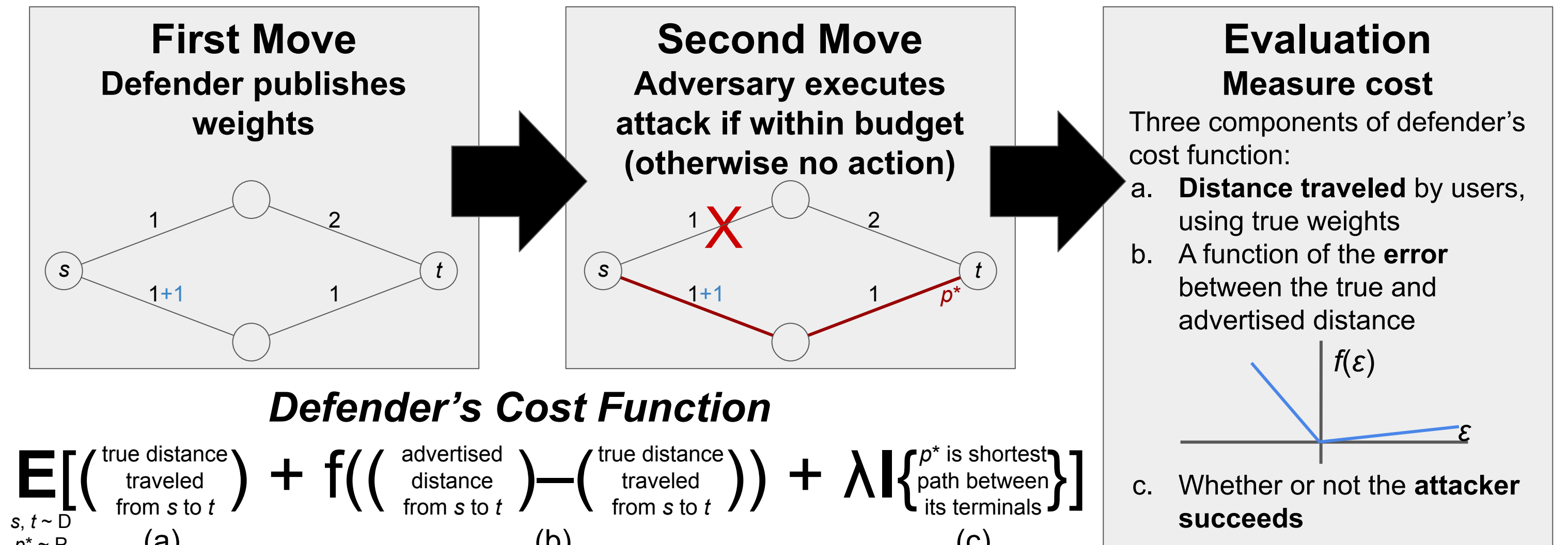
## Results

- Graphs with weights drawn from a Poisson distribution
- Edge removal costs are equal
- Randomly selected  $s$  and  $t$
- Defender uncertainty: distribution of target paths and adversary budgets
- Users are more likely to travel between nodes on potential target paths

## Optimization Approach

<https://arxiv.org/abs/2305.19083>

### Stackelberg Game



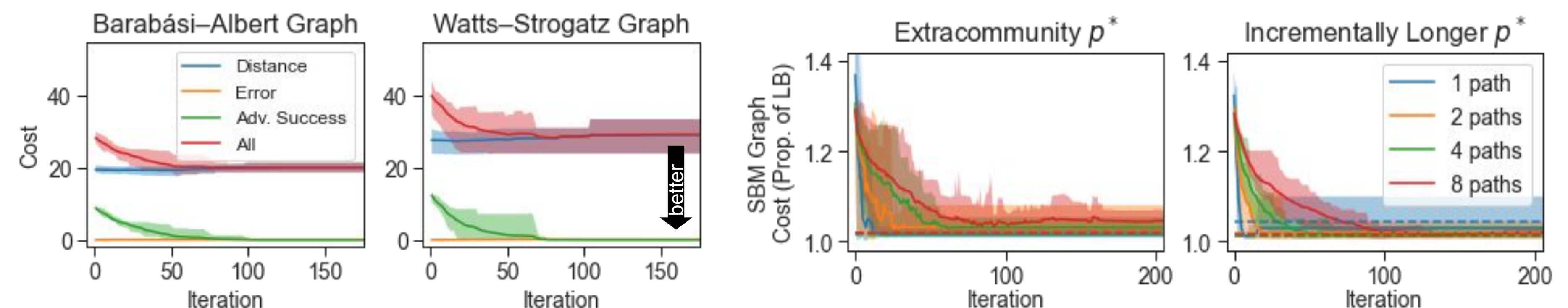
### Defender's Cost Function

$$E\left[\left(\frac{\text{true distance traveled from } s \text{ to } t}{s, t \sim D}\right) + f\left(\left(\frac{\text{advertised distance traveled from } s \text{ to } t}{p^* \sim P}\right) - \left(\frac{\text{true distance traveled from } s \text{ to } t}{b \sim B}\right)\right) + \lambda I\{p^* \text{ is shortest path between its terminals}\}\right]$$

- Cut Defense problem: Optimize defender's cost
- Theorem: Zero-sum Cut Defense with an attack oracle is NP-hard

### Algorithm: PATHDEFENSE

- Heuristic algorithm to optimize defender cost
- Iteratively increment edges on  $p^*$  to decrease probability of attack
- After termination: Choose lowest-cost weights



- Attack probability becomes negligible with small increase in distance
- Cost due to error is extremely low

Defense is more difficult and costly if it exits and returns to the community of  $s$  and  $t$